

Theoretical Computer Science Cheat Sheet

| Definitions | | Series | |
|---|--|---|--|
| $f(n) = O(g(n))$ | iff \exists positive c, n_0 such that $0 \leq f(n) \leq cg(n) \forall n \geq n_0$. | $\sum_{i=1}^n i = \frac{n(n+1)}{2}, \quad \sum_{i=1}^n i^2 = \frac{n(n+1)(2n+1)}{6}, \quad \sum_{i=1}^n i^3 = \frac{n^2(n+1)^2}{4}$ | |
| $f(n) = \Omega(g(n))$ | iff \exists positive c, n_0 such that $f(n) \geq cg(n) \geq 0 \forall n \geq n_0$. | In general: $\sum_{i=1}^n i^m = \frac{1}{m+1} \left[(n+1)^{m+1} - 1 - \sum_{i=1}^n ((i+1)^{m+1} - i^{m+1} - (m+1)i^m) \right]$ $\sum_{i=1}^{n-1} i^m = \frac{1}{m+1} \sum_{k=0}^m \binom{m+1}{k} B_k n^{m+1-k}.$ | |
| $f(n) = \Theta(g(n))$ | iff $f(n) = O(g(n))$ and $f(n) = \Omega(g(n))$. | | |
| $f(n) = o(g(n))$ | iff $\lim_{n \rightarrow \infty} f(n)/g(n) = 0$. | | |
| $\lim_{n \rightarrow \infty} a_n = a$ | iff $\forall \epsilon > 0, \exists n_0$ such that $ a_n - a < \epsilon, \forall n \geq n_0$. | | |
| $\sup S$ | least $b \in \mathbb{R}$ such that $b \geq s, \forall s \in S$. | | |
| $\inf S$ | greatest $b \in \mathbb{R}$ such that $b \leq s, \forall s \in S$. | Geometric series: $\sum_{i=0}^n c^i = \frac{c^{n+1} - 1}{c - 1}, \quad c \neq 1, \quad \sum_{i=0}^{\infty} c^i = \frac{1}{1 - c}, \quad \sum_{i=1}^{\infty} c^i = \frac{c}{1 - c}, \quad c < 1,$ $\sum_{i=0}^n i c^i = \frac{nc^{n+2} - (n+1)c^{n+1} + c}{(c-1)^2}, \quad c \neq 1, \quad \sum_{i=0}^{\infty} i c^i = \frac{c}{(1-c)^2}, \quad c < 1.$ | |
| $\liminf_{n \rightarrow \infty} a_n$ | $\liminf_{n \rightarrow \infty} \{a_i \mid i \geq n, i \in \mathbb{N}\}$. | Harmonic series: $H_n = \sum_{i=1}^n \frac{1}{i}, \quad \sum_{i=1}^n i H_i = \frac{n(n+1)}{2} H_n - \frac{n(n-1)}{4}.$ $\sum_{i=1}^n H_i = (n+1)H_n - n, \quad \sum_{i=1}^n \binom{i}{m} H_i = \binom{n+1}{m+1} \left(H_{n+1} - \frac{1}{m+1} \right).$ | |
| $\binom{n}{k}$ | Combinations: Size k subsets of a size n set. | 1. $\binom{n}{k} = \frac{n!}{(n-k)!k!}, \quad 2. \sum_{k=0}^n \binom{n}{k} = 2^n, \quad 3. \binom{n}{k} = \binom{n}{n-k},$ 4. $\binom{n}{k} = \frac{n}{k} \binom{n-1}{k-1}, \quad 5. \binom{n}{k} = \binom{n-1}{k} + \binom{n-1}{k-1},$ 6. $\binom{n}{m} \binom{m}{k} = \binom{n}{k} \binom{n-k}{m-k}, \quad 7. \sum_{k=0}^n \binom{r+k}{k} = \binom{r+n+1}{n},$ 8. $\sum_{k=0}^n \binom{k}{m} = \binom{n+1}{m+1}, \quad 9. \sum_{k=0}^n \binom{r}{k} \binom{s}{n-k} = \binom{r+s}{n},$ 10. $\binom{n}{k} = (-1)^k \binom{k-n-1}{k}, \quad 11. \begin{Bmatrix} n \\ 1 \end{Bmatrix} = \begin{Bmatrix} n \\ n \end{Bmatrix} = 1,$ 12. $\begin{Bmatrix} n \\ 2 \end{Bmatrix} = 2^{n-1} - 1, \quad 13. \begin{Bmatrix} n \\ k \end{Bmatrix} = k \begin{Bmatrix} n-1 \\ k \end{Bmatrix} + \begin{Bmatrix} n-1 \\ k-1 \end{Bmatrix},$ | |
| $\begin{Bmatrix} n \\ k \end{Bmatrix}$ | Stirling numbers (1st kind): Arrangements of an n element set into k cycles. | | |
| $\begin{Bmatrix} n \\ k \end{Bmatrix}$ | Stirling numbers (2nd kind): Partitions of an n element set into k non-empty sets. | | |
| $\langle \begin{smallmatrix} n \\ k \end{smallmatrix} \rangle$ | 1st order Eulerian numbers: Permutations $\pi_1 \pi_2 \dots \pi_n$ on $\{1, 2, \dots, n\}$ with k ascents. | | |
| $\langle\!\langle \begin{smallmatrix} n \\ k \end{smallmatrix} \rangle\!\rangle$ | 2nd order Eulerian numbers. | | |
| C_n | Catalan Numbers: Binary trees with $n+1$ vertices. | | |
| 14. $\begin{Bmatrix} n \\ 1 \end{Bmatrix} = (n-1)!$ | 15. $\begin{Bmatrix} n \\ 2 \end{Bmatrix} = (n-1)!H_{n-1}$ | 16. $\begin{Bmatrix} n \\ n \end{Bmatrix} = 1$ | 17. $\begin{Bmatrix} n \\ k \end{Bmatrix} \geq \begin{Bmatrix} n \\ k \end{Bmatrix}$ |
| 18. $\begin{Bmatrix} n \\ k \end{Bmatrix} = (n-1) \begin{Bmatrix} n-1 \\ k \end{Bmatrix} + \begin{Bmatrix} n-1 \\ k-1 \end{Bmatrix}$ | 19. $\begin{Bmatrix} n \\ n-1 \end{Bmatrix} = \begin{Bmatrix} n \\ n-1 \end{Bmatrix} = \binom{n}{2}$ | 20. $\sum_{k=0}^n \begin{Bmatrix} n \\ k \end{Bmatrix} = n!$ | 21. $C_n = \frac{1}{n+1} \binom{2n}{n}$ |
| 22. $\begin{Bmatrix} n \\ 0 \end{Bmatrix} = \begin{Bmatrix} n \\ n-1 \end{Bmatrix} = 1$ | 23. $\begin{Bmatrix} n \\ k \end{Bmatrix} = \begin{Bmatrix} n \\ n-1-k \end{Bmatrix}$ | 24. $\begin{Bmatrix} n \\ k \end{Bmatrix} = (k+1) \begin{Bmatrix} n-1 \\ k \end{Bmatrix} + (n-k) \begin{Bmatrix} n-1 \\ k-1 \end{Bmatrix}$ | |
| 25. $\begin{Bmatrix} 0 \\ k \end{Bmatrix} = \begin{cases} 1 & \text{if } k=0, \\ 0 & \text{otherwise} \end{cases}$ | 26. $\begin{Bmatrix} n \\ 1 \end{Bmatrix} = 2^n - n - 1$ | 27. $\begin{Bmatrix} n \\ 2 \end{Bmatrix} = 3^n - (n+1)2^n + \binom{n+1}{2}$ | |
| 28. $x^n = \sum_{k=0}^n \begin{Bmatrix} n \\ k \end{Bmatrix} \binom{x+k}{n}$ | 29. $\begin{Bmatrix} n \\ m \end{Bmatrix} = \sum_{k=0}^m \binom{n+1}{k} (m+1-k)^n (-1)^k$ | 30. $m! \begin{Bmatrix} n \\ m \end{Bmatrix} = \sum_{k=0}^n \begin{Bmatrix} n \\ k \end{Bmatrix} \binom{k}{n-m}$ | |
| 31. $\begin{Bmatrix} n \\ m \end{Bmatrix} = \sum_{k=0}^n \begin{Bmatrix} n \\ k \end{Bmatrix} \binom{n-k}{m} (-1)^{n-k-m} k!$ | 32. $\langle\!\langle\!\langle \begin{smallmatrix} n \\ 0 \end{smallmatrix} \rangle\!\rangle\!\rangle = 1$ | 33. $\langle\!\langle\!\langle \begin{smallmatrix} n \\ n \end{smallmatrix} \rangle\!\rangle\!\rangle = 0 \quad \text{for } n \neq 0$ | |
| 34. $\langle\!\langle\!\langle \begin{smallmatrix} n \\ k \end{smallmatrix} \rangle\!\rangle\!\rangle = (k+1) \langle\!\langle\!\langle \begin{smallmatrix} n-1 \\ k \end{smallmatrix} \rangle\!\rangle\!\rangle + (2n-1-k) \langle\!\langle\!\langle \begin{smallmatrix} n-1 \\ k-1 \end{smallmatrix} \rangle\!\rangle\!\rangle$ | | 35. $\sum_{k=0}^n \langle\!\langle\!\langle \begin{smallmatrix} n \\ k \end{smallmatrix} \rangle\!\rangle\!\rangle = \frac{(2n)n}{2^n}$ | |

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Identities Cont.

| | |
|--|--|
| 38. $\begin{bmatrix} n+1 \\ m+1 \end{bmatrix} = \sum_k \begin{bmatrix} n \\ k \end{bmatrix} \binom{k}{m} = \sum_{k=0}^n \begin{bmatrix} k \\ m \end{bmatrix} n^{n-k} = n! \sum_{k=0}^n \frac{1}{k!} \begin{bmatrix} k \\ m \end{bmatrix}$, 40. $\begin{Bmatrix} n \\ m \end{Bmatrix} = \sum_k \binom{n}{k} \begin{Bmatrix} k+1 \\ m+1 \end{Bmatrix} (-1)^{n-k}$, 42. $\begin{Bmatrix} m+n+1 \\ m \end{Bmatrix} = \sum_{k=0}^m k \begin{Bmatrix} n+k \\ k \end{Bmatrix}$, 44. $\binom{n}{m} = \sum_k \begin{Bmatrix} n+1 \\ k+1 \end{Bmatrix} \binom{k}{m} (-1)^{m-k}$, 46. $\begin{Bmatrix} n \\ n-m \end{Bmatrix} = \sum_k \binom{m-n}{m+k} \binom{m+n}{n+k} \binom{m+k}{k}$, 48. $\begin{Bmatrix} n \\ \ell+m \end{Bmatrix} \binom{\ell+m}{\ell} = \sum_k \begin{Bmatrix} k \\ \ell \end{Bmatrix} \begin{Bmatrix} n-k \\ m \end{Bmatrix} \binom{n}{k}$, | 39. $\begin{bmatrix} x \\ x-n \end{bmatrix} = \sum_{k=0}^n \begin{Bmatrix} n \\ k \end{Bmatrix} \binom{x+k}{2n}$, 41. $\begin{bmatrix} n \\ m \end{bmatrix} = \sum_k \begin{bmatrix} n+1 \\ k+1 \end{Bmatrix} \binom{k}{m} (-1)^{m-k}$, 43. $\begin{bmatrix} m+n+1 \\ m \end{bmatrix} = \sum_{k=0}^m k(n+k) \binom{n+k}{k}$, 45. $(n-m)! \binom{n}{m} = \sum_k \begin{Bmatrix} n+1 \\ k+1 \end{Bmatrix} \binom{k}{m} (-1)^{m-k}$, for $n \geq m$, 47. $\begin{bmatrix} n \\ n-m \end{Bmatrix} = \sum_k \binom{m-n}{m+k} \binom{m+n}{n+k} \binom{m+k}{k}$, 49. $\begin{bmatrix} n \\ \ell+m \end{Bmatrix} \binom{\ell+m}{\ell} = \sum_k \begin{Bmatrix} k \\ \ell \end{Bmatrix} \binom{n-k}{m} \binom{n}{k}$. |
|--|--|

Every vertex has at most two edges.
 Kraft inequality: If the depths of the leaves of a binary tree are d_1, \dots, d_n :

$$\sum_{i=1}^n 2^{-d_i} \leq 1$$

 and equality holds only if every internal node has 2 sons.

Recurrences

Master method:

$$T(n) = aT(n/b) + f(n), \quad a \geq 1, b > 1$$

If $\exists \epsilon > 0$ such that $f(n) = O(n^{\log_b a - \epsilon})$ then

$$T(n) = \Theta(n^{\log_b a}).$$

If $f(n) = \Theta(n^{\log_b a})$ then

$$T(n) = \Theta(n^{\log_b a} \log_2 n).$$

If $\exists \epsilon > 0$ such that $f(n) = \Omega(n^{\log_b a + \epsilon})$, and $\exists c < 1$ such that $aT(n/b) \leq cf(n)$ for large n , then

$$T(n) = \Theta(f(n)).$$

Substitution (example): Consider the following recurrence

$$T_{i+1} = 2^{2^i} \cdot T_i^2, \quad T_1 = 2.$$

Note that T_i is always a power of two.

Let $t_i = \log_2 T_i$. Then we have

$$t_{i+1} = 2^i + 2t_i, \quad t_1 = 1.$$

Let $u_i = t_i/2^i$. Dividing both sides of the previous equation by 2^{i+1} we get

$$\frac{t_{i+1}}{2^{i+1}} = \frac{2^i}{2^{i+1}} + \frac{t_i}{2^i}.$$

Substituting we find

$$u_{i+1} = \frac{1}{2} + u_i, \quad u_1 = \frac{1}{2},$$

which is simply $u_i = i/2$. So we find that T_i has the closed form $T_i = 2^{i2^{i-1}}$.

Summing factors (example): Consider the following recurrence

$$T(n) = 3T(n/2) + n, \quad T(1) = 1.$$

Rewrite so that all terms involving T are on the left side

$$T(n) - 3T(n/2) = n.$$

$$\begin{aligned} 1(T(n) - 3T(n/2)) &= n \\ 3(T(n/2) - 3T(n/4)) &= n/2 \\ &\vdots & \vdots & \vdots \\ 3^{\log_2 n-1}(T(2) - 3T(1)) &= 2 \end{aligned}$$

Let $m = \log_2 n$. Summing the left side we get $T(n) - 3^m T(1) = T(n) - 3^m = T(n) - n^k$ where $k = \log_2 3 \approx 1.58496$. Summing the right side we get

$$\sum_{i=0}^{m-1} \frac{n}{2^i} 3^i = n \sum_{i=0}^{m-1} \left(\frac{3}{2}\right)^i.$$

Let $c = \frac{3}{2}$. Then we have

$$\begin{aligned} n \sum_{i=0}^{m-1} c^i &= n \left(\frac{c^m - 1}{c - 1} \right) \\ &= 2n(c^{\log_2 n} - 1) \\ &= 2n(c^{(k-1)\log_c n} - 1) \\ &= 2n^k - 2n, \end{aligned}$$

and so $T(n) = 3n^k - 2n$. Full history recurrences can often be changed to limited history ones (example): Consider

$$T_i = 1 + \sum_{j=0}^{i-1} T_j, \quad T_0 = 1.$$

Note that

$$T_{i+1} = 1 + \sum_{j=0}^i T_j.$$

Subtracting we find

$$T_{i+1} - T_i = 1 + \sum_{j=0}^i T_j - 1 - \sum_{j=0}^{i-1} T_j$$

Generating functions:

- Multiply both sides of the equation by x^i .
 - Sum both sides over all i for which the equation is valid.
 - Choose a generating function $G(x)$. Usually $G(x) = \sum_{i=0}^{\infty} x^i g_i$.
 - Rewrite the equation in terms of the generating function $G(x)$.
 - Solve for $G(x)$.
 - The coefficient of x^i in $G(x)$ is g_i .
- Example:

$$g_{i+1} = 2g_i + 1, \quad g_0 = 0.$$

Multiply and sum:

$$\sum_{i \geq 0} g_{i+1} x^i = \sum_{i \geq 0} 2g_i x^i + \sum_{i \geq 0} x^i.$$

We choose $G(x) = \sum_{i \geq 0} x^i g_i$. Rewrite in terms of $G(x)$:

$$\frac{G(x) - g_0}{x} = 2G(x) + \sum_{i \geq 0} x^i.$$

Simplify:

$$\frac{G(x)}{x} = 2G(x) + \frac{1}{1-x}.$$

Solve for $G(x)$:

$$G(x) = \frac{x}{(1-x)(1-2x)}.$$

Expand this using partial fractions:

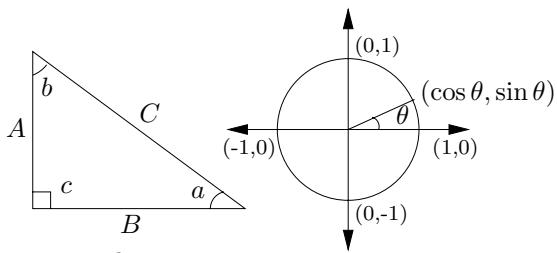
$$\begin{aligned} G(x) &= x \left(\frac{2}{1-2x} - \frac{1}{1-x} \right) \\ &= x \left(2 \sum_{i \geq 0} 2^i x^i - \sum_{i \geq 0} x^i \right) \\ &= \sum (2^{i+1} - 1)x^{i+1}. \end{aligned}$$

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| $\pi \approx 3.14159$, | $e \approx 2.71828$, | $\gamma \approx 0.57721$, | $\phi = \frac{1+\sqrt{5}}{2} \approx 1.61803$, | $\hat{\phi} = \frac{1-\sqrt{5}}{2} \approx -0.61803$ |
|---|-----------------------|----------------------------|--|---|
| i | 2^i | p_i | General | Probability |
| 1 | 2 | 2 | Bernoulli Numbers ($B_i = 0$, odd $i \neq 1$): $B_0 = 1$, $B_1 = -\frac{1}{2}$, $B_2 = \frac{1}{4}$, $B_4 = -\frac{1}{30}$, $B_6 = \frac{1}{42}$, $B_8 = -\frac{1}{30}$, $B_{10} = \frac{5}{66}$. | Continuous distributions: If $\Pr[a < X < b] = \int_a^b p(x) dx$, then p is the probability density function of X . If $\Pr[X < a] = P(a)$, |
| 2 | 4 | 3 | Change of base, quadratic formula: $\log_b x = \frac{\log_a x}{\log_a b}, \quad \frac{-b \pm \sqrt{b^2 - 4ac}}{2a}$. | then P is the distribution function of X . If P and p both exist then $P(a) = \int_{-\infty}^a p(x) dx$. |
| 3 | 8 | 5 | Euler's number e : $e = 1 + \frac{1}{2} + \frac{1}{6} + \frac{1}{24} + \frac{1}{120} + \dots$ $\lim_{n \rightarrow \infty} \left(1 + \frac{x}{n}\right)^n = e^x$. | Expectation: If X is discrete $E[g(X)] = \sum_x g(x) \Pr[X = x]$. If X continuous then $E[g(X)] = \int_{-\infty}^{\infty} g(x)p(x) dx = \int_{-\infty}^{\infty} g(x) dP(x)$. |
| 4 | 16 | 7 | $(1 + \frac{1}{n})^n < e < (1 + \frac{1}{n})^{n+1}$. | Variance, standard deviation: $\text{VAR}[X] = E[X^2] - E[X]^2$, $\sigma = \sqrt{\text{VAR}[X]}$. |
| 5 | 32 | 11 | $(1 + \frac{1}{n})^n = e - \frac{e}{2n} + \frac{11e}{24n^2} - O\left(\frac{1}{n^3}\right)$. | For events A and B : $\Pr[A \vee B] = \Pr[A] + \Pr[B] - \Pr[A \wedge B]$ |
| 6 | 64 | 13 | Harmonic numbers: $1, \frac{3}{2}, \frac{11}{6}, \frac{25}{12}, \frac{137}{60}, \frac{49}{20}, \frac{363}{140}, \frac{761}{280}, \frac{7129}{2520}, \dots$ | $\Pr[A \wedge B] = \Pr[A] \cdot \Pr[B]$, iff A and B are independent. $\Pr[A B] = \frac{\Pr[A \wedge B]}{\Pr[B]}$ |
| 7 | 128 | 17 | $\ln n < H_n < \ln n + 1$, $H_n = \ln n + \gamma + O\left(\frac{1}{n}\right)$. | For random variables X and Y : $E[X \cdot Y] = E[X] \cdot E[Y]$, if X and Y are independent. |
| 8 | 256 | 19 | Factorial, Stirling's approximation: $1, 2, 6, 24, 120, 720, 5040, 40320, 362880, \dots$ | $E[X + Y] = E[X] + E[Y]$, $E[cX] = cE[X]$. |
| 9 | 512 | 23 | $n! = \sqrt{2\pi n} \left(\frac{n}{e}\right)^n \left(1 + \Theta\left(\frac{1}{n}\right)\right)$. | Bayes' theorem: $\Pr[A_i B] = \frac{\Pr[B A_i] \Pr[A_i]}{\sum_{j=1}^n \Pr[A_j] \Pr[B A_j]}$. |
| 10 | 1,024 | 29 | Ackermann's function and inverse: $a(i, j) = \begin{cases} 2^j & i = 1 \\ a(i-1, 2) & j = 1 \\ a(i-1, a(i, j-1)) & i, j \geq 2 \end{cases}$ | Inclusion-exclusion: $\Pr\left[\bigvee_{i=1}^n X_i\right] = \sum_{i=1}^n \Pr[X_i] + \sum_{k=2}^n (-1)^{k+1} \sum_{i_1 < \dots < i_k} \Pr\left[\bigwedge_{j=1}^k X_{i_j}\right]$. |
| 11 | 2,048 | 31 | $\alpha(i) = \min\{j \mid a(j, j) \geq i\}$. | Moment inequalities: $\Pr[X \geq \lambda E[X]] \leq \frac{1}{\lambda}$, $\Pr[X - E[X] \geq \lambda \cdot \sigma] \leq \frac{1}{\lambda^2}$. |
| 12 | 4,096 | 37 | Binomial distribution: $\Pr[X = k] = \binom{n}{k} p^k q^{n-k}, \quad q = 1 - p$, $E[X] = \sum_{k=1}^n k \binom{n}{k} p^k q^{n-k} = np$. | Geometric distribution: $\Pr[X = k] = pq^{k-1}, \quad q = 1 - p$, |
| 13 | 8,192 | 41 | Poisson distribution: $\Pr[X = k] = \frac{e^{-\lambda} \lambda^k}{k!}, \quad E[X] = \lambda$. | |
| 14 | 16,384 | 43 | Normal (Gaussian) distribution: $p(x) = \frac{1}{\sqrt{2\pi\sigma^2}} e^{-(x-\mu)^2/2\sigma^2}, \quad E[X] = \mu$. | |
| 15 | 32,768 | 47 | The "coupon collector": We are given a random coupon each day, and there are n different types of coupons. The distribution of coupons is uniform. The expected number of days to pass before we collect all n coupons is $E[X] = n \cdot H_n$. | |
| 16 | 65,536 | 53 | | |
| 17 | 131,072 | 59 | | |
| 18 | 262,144 | 61 | | |
| 19 | 524,288 | 67 | | |
| 20 | 1,048,576 | 71 | | |
| 21 | 2,097,152 | 73 | | |
| 22 | 4,194,304 | 79 | | |
| 23 | 8,388,608 | 83 | | |
| 24 | 16,777,216 | 89 | | |
| 25 | 33,554,432 | 97 | | |
| 26 | 67,108,864 | 101 | | |
| 27 | 134,217,728 | 103 | | |
| 28 | 268,435,456 | 107 | | |
| 29 | 536,870,912 | 109 | | |
| 30 | 1,073,741,824 | 113 | | |
| 31 | 2,147,483,648 | 127 | | |
| 32 | 4,294,967,296 | 131 | | |
| Pascal's Triangle | | | | |
| 1 1 1 1 2 1 1 3 3 1 1 4 6 4 1 1 5 10 10 5 1 1 6 15 20 15 6 1 1 7 21 35 35 21 7 1 | | | | |

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Trigonometry



Pythagorean theorem:

$$C^2 = A^2 + B^2.$$

Definitions:

$$\sin a = A/C, \quad \cos a = B/C,$$

$$\csc a = C/A, \quad \sec a = C/B,$$

$$\tan a = \frac{\sin a}{\cos a} = \frac{A}{B}, \quad \cot a = \frac{\cos a}{\sin a} = \frac{B}{A}.$$

Area, radius of inscribed circle:

$$\frac{1}{2}AB, \quad \frac{AB}{A+B+C}.$$

Identities:

$$\sin x = \frac{1}{\csc x}, \quad \cos x = \frac{1}{\sec x},$$

$$\tan x = \frac{1}{\cot x}, \quad \sin^2 x + \cos^2 x = 1,$$

$$1 + \tan^2 x = \sec^2 x, \quad 1 + \cot^2 x = \csc^2 x,$$

$$\sin x = \cos(\frac{\pi}{2} - x), \quad \sin x = \sin(\pi - x),$$

$$\cos x = -\cos(\pi - x), \quad \tan x = \cot(\frac{\pi}{2} - x),$$

$$\cot x = -\cot(\pi - x), \quad \csc x = \cot \frac{x}{2} - \cot x,$$

$$\sin(x \pm y) = \sin x \cos y \pm \cos x \sin y,$$

$$\cos(x \pm y) = \cos x \cos y \mp \sin x \sin y,$$

$$\tan(x \pm y) = \frac{\tan x \pm \tan y}{1 \mp \tan x \tan y},$$

$$\cot(x \pm y) = \frac{\cot x \cot y \mp 1}{\cot x \pm \cot y},$$

$$\sin 2x = 2 \sin x \cos x, \quad \sin 2x = \frac{2 \tan x}{1 + \tan^2 x},$$

$$\cos 2x = \cos^2 x - \sin^2 x, \quad \cos 2x = 2 \cos^2 x - 1,$$

$$\cos 2x = 1 - 2 \sin^2 x, \quad \cos 2x = \frac{1 - \tan^2 x}{1 + \tan^2 x},$$

$$\tan 2x = \frac{2 \tan x}{1 - \tan^2 x}, \quad \cot 2x = \frac{\cot^2 x - 1}{2 \cot x},$$

$$\sin(x+y) \sin(x-y) = \sin^2 x - \sin^2 y,$$

$$\cos(x+y) \cos(x-y) = \cos^2 x - \sin^2 y.$$

$$\text{Euler's equation: } e^{ix} = \cos x + i \sin x, \quad e^{i\pi} = -1.$$

Matrices

Multiplication:

$$C = A \cdot B, \quad c_{i,j} = \sum_{k=1}^n a_{i,k} b_{k,j}.$$

Determinants: $\det A \neq 0$ iff A is non-singular.

$$\det A \cdot B = \det A \cdot \det B,$$

$$\det A = \sum_{\pi} \prod_{i=1}^n \text{sign}(\pi) a_{i,\pi(i)}.$$

2×2 and 3×3 determinant:

$$\begin{vmatrix} a & b \\ c & d \end{vmatrix} = ad - bc,$$

$$\begin{vmatrix} a & b & c \\ d & e & f \\ g & h & i \end{vmatrix} = g \begin{vmatrix} b & c \\ e & f \end{vmatrix} - h \begin{vmatrix} a & c \\ d & f \end{vmatrix} + i \begin{vmatrix} a & b \\ d & e \end{vmatrix}$$

$$= aei + bfg + cdh - ceg - fha - ibd.$$

Permanents:

$$\text{perm } A = \sum_{\pi} \prod_{i=1}^n a_{i,\pi(i)}.$$

Hyperbolic Functions

Definitions:

$$\sinh x = \frac{e^x - e^{-x}}{2}, \quad \cosh x = \frac{e^x + e^{-x}}{2},$$

$$\tanh x = \frac{e^x - e^{-x}}{e^x + e^{-x}}, \quad \csch x = \frac{1}{\sinh x},$$

$$\sech x = \frac{1}{\cosh x}, \quad \coth x = \frac{1}{\tanh x}.$$

Identities:

$$\cosh^2 x - \sinh^2 x = 1, \quad \tanh^2 x + \sech^2 x = 1,$$

$$\coth^2 x - \csch^2 x = 1, \quad \sinh(-x) = -\sinh x,$$

$$\cosh(-x) = \cosh x, \quad \tanh(-x) = -\tanh x,$$

$$\sinh(x+y) = \sinh x \cosh y + \cosh x \sinh y,$$

$$\cosh(x+y) = \cosh x \cosh y + \sinh x \sinh y,$$

$$\sinh 2x = 2 \sinh x \cosh x,$$

$$\cosh 2x = \cosh^2 x + \sinh^2 x,$$

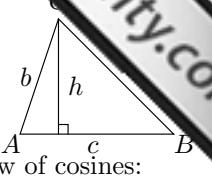
$$\cosh x + \sinh x = e^x, \quad \cosh x - \sinh x = e^{-x},$$

$$(\cosh x + \sinh x)^n = \cosh nx + \sinh nx, \quad n \in \mathbb{Z},$$

$$2 \sinh^2 \frac{x}{2} = \cosh x - 1, \quad 2 \cosh^2 \frac{x}{2} = \cosh x + 1.$$

| θ | $\sin \theta$ | $\cos \theta$ | $\tan \theta$ |
|-----------------|----------------------|----------------------|----------------------|
| 0 | 0 | 1 | 0 |
| $\frac{\pi}{6}$ | $\frac{1}{2}$ | $\frac{\sqrt{3}}{2}$ | $\frac{\sqrt{3}}{3}$ |
| $\frac{\pi}{4}$ | $\frac{\sqrt{2}}{2}$ | $\frac{\sqrt{2}}{2}$ | 1 |

... in mathematics you don't understand things, you just get used to them.



Law of cosines:

$$c^2 = a^2 + b^2 - 2ab \cos C.$$

Area:

$$A = \frac{1}{2}hc,$$

$$= \frac{1}{2}ab \sin C,$$

$$= \frac{c^2 \sin A \sin B}{2 \sin C}.$$

Heron's formula:

$$A = \sqrt{s \cdot s_a \cdot s_b \cdot s_c},$$

$$s = \frac{1}{2}(a+b+c),$$

$$s_a = s - a,$$

$$s_b = s - b,$$

$$s_c = s - c.$$

More identities:

$$\sin \frac{x}{2} = \sqrt{\frac{1 - \cos x}{2}},$$

$$\cos \frac{x}{2} = \sqrt{\frac{1 + \cos x}{2}},$$

$$\tan \frac{x}{2} = \sqrt{\frac{1 - \cos x}{1 + \cos x}},$$

$$= \frac{1 - \cos x}{\sin x},$$

$$= \frac{\sin x}{1 + \cos x},$$

$$\cot \frac{x}{2} = \sqrt{\frac{1 + \cos x}{1 - \cos x}},$$

$$= \frac{1 + \cos x}{\sin x},$$

$$= \frac{\sin x}{1 - \cos x},$$

$$\sin x = \frac{e^{ix} - e^{-ix}}{2i},$$

$$\cos x = \frac{e^{ix} + e^{-ix}}{2},$$

$$\tan x = -i \frac{e^{ix} - e^{-ix}}{e^{ix} + e^{-ix}},$$

$$= -i \frac{e^{2ix} - 1}{e^{2ix} + 1},$$

$$\sinh x = \frac{e^{ix} - e^{-ix}}{i},$$

$$\cosh x = \cosh ix,$$

+ sinh ix

Theoretical Computer Science Cheat Sheet

| Number Theory | Graph Theory | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | |
|--|--|-------------|--|-----------------|----------------------------|---------------|-------------------------------------|-------------|--|--------------|-----------------------------|-------------|---------------------------------|------------------|---|------------------|-------------------------------|-------------|----------------------------|------------------|----------------------|------------|-------------------------|-----------------|---|--------------------|---|------------|--|----------------|----------------|-----------------|---------------|--------------------|---|----------------|---|------------------|--|-----------------|-----------------------------------|-----------------|---|---------------|---|-----------------|--|---------------------|--|---------------------|---|--------------------|---------------------------------|
| <p>The Chinese remainder theorem: There exists a number C such that:</p> $C \equiv r_1 \pmod{m_1}$ $\vdots \quad \vdots \quad \vdots$ $C \equiv r_n \pmod{m_n}$ <p>if m_i and m_j are relatively prime for $i \neq j$.</p> <p>Euler's function: $\phi(x)$ is the number of positive integers less than x relatively prime to x. If $\prod_{i=1}^n p_i^{e_i}$ is the prime factorization of x then</p> $\phi(x) = \prod_{i=1}^n p_i^{e_i-1} (p_i - 1).$ <p>Euler's theorem: If a and b are relatively prime then</p> $1 \equiv a^{\phi(b)} \pmod{b}.$ <p>Fermat's theorem:</p> $1 \equiv a^{p-1} \pmod{p}.$ <p>The Euclidean algorithm: if $a > b$ are integers then</p> $\gcd(a, b) = \gcd(a \bmod b, b).$ <p>If $\prod_{i=1}^n p_i^{e_i}$ is the prime factorization of x then</p> $S(x) = \sum_{d x} d = \prod_{i=1}^n \frac{p_i^{e_i+1} - 1}{p_i - 1}.$ <p>Perfect Numbers: x is an even perfect number iff $x = 2^{n-1}(2^n - 1)$ and $2^n - 1$ is prime.</p> <p>Wilson's theorem: n is a prime iff</p> $(n-1)! \equiv -1 \pmod{n}.$ <p>Möbius inversion:</p> $\mu(i) = \begin{cases} 1 & \text{if } i = 1. \\ 0 & \text{if } i \text{ is not square-free.} \\ (-1)^r & \text{if } i \text{ is the product of } r \text{ distinct primes.} \end{cases}$ <p>If</p> $G(a) = \sum_{d a} F(d),$ <p>then</p> $F(a) = \sum_{d a} \mu(d) G\left(\frac{a}{d}\right).$ <p>Prime numbers:</p> $p_n = n \ln n + n \ln \ln n - n + n \frac{\ln \ln n}{\ln n} + O\left(\frac{n}{\ln n}\right),$ $\pi(n) = \frac{n}{\ln n} + \frac{n}{(\ln n)^2} + \frac{2!n}{(\ln n)^3}$ | <p>Definitions:</p> <table> <tr> <td><i>Loop</i></td><td>An edge connecting a vertex to itself.</td></tr> <tr> <td><i>Directed</i></td><td>Each edge has a direction.</td></tr> <tr> <td><i>Simple</i></td><td>Graph with no loops or multi-edges.</td></tr> <tr> <td><i>Walk</i></td><td>A sequence $v_0 e_1 v_1 \dots e_\ell v_\ell$.</td></tr> <tr> <td><i>Trail</i></td><td>A walk with distinct edges.</td></tr> <tr> <td><i>Path</i></td><td>A trail with distinct vertices.</td></tr> <tr> <td><i>Connected</i></td><td>A graph where there exists a path between any two vertices.</td></tr> <tr> <td><i>Component</i></td><td>A maximal connected subgraph.</td></tr> <tr> <td><i>Tree</i></td><td>A connected acyclic graph.</td></tr> <tr> <td><i>Free tree</i></td><td>A tree with no root.</td></tr> <tr> <td><i>DAG</i></td><td>Directed acyclic graph.</td></tr> <tr> <td><i>Eulerian</i></td><td>Graph with a trail visiting each edge exactly once.</td></tr> <tr> <td><i>Hamiltonian</i></td><td>Graph with a cycle visiting each vertex exactly once.</td></tr> <tr> <td><i>Cut</i></td><td>A set of edges whose removal increases the number of components.</td></tr> <tr> <td><i>Cut-set</i></td><td>A minimal cut.</td></tr> <tr> <td><i>Cut edge</i></td><td>A size 1 cut.</td></tr> <tr> <td><i>k-Connected</i></td><td>A graph connected with the removal of any $k-1$ vertices.</td></tr> <tr> <td><i>k-Tough</i></td><td>$\forall S \subseteq V, S \neq \emptyset$ we have $k \cdot c(G-S) \leq S$.</td></tr> <tr> <td><i>k-Regular</i></td><td>A graph where all vertices have degree k.</td></tr> <tr> <td><i>k-Factor</i></td><td>A k-regular spanning subgraph.</td></tr> <tr> <td><i>Matching</i></td><td>A set of edges, no two of which are adjacent.</td></tr> <tr> <td><i>Clique</i></td><td>A set of vertices, all of which are adjacent.</td></tr> <tr> <td><i>Ind. set</i></td><td>A set of vertices, none of which are adjacent.</td></tr> <tr> <td><i>Vertex cover</i></td><td>A set of vertices which cover all edges.</td></tr> <tr> <td><i>Planar graph</i></td><td>A graph which can be embedded in the plane.</td></tr> <tr> <td><i>Plane graph</i></td><td>An embedding of a planar graph.</td></tr> </table> <p style="text-align: center;">$\sum_{v \in V} \deg(v) = 2m.$</p> <p>If G is planar then $n - m + f = 2$, so</p> | <i>Loop</i> | An edge connecting a vertex to itself. | <i>Directed</i> | Each edge has a direction. | <i>Simple</i> | Graph with no loops or multi-edges. | <i>Walk</i> | A sequence $v_0 e_1 v_1 \dots e_\ell v_\ell$. | <i>Trail</i> | A walk with distinct edges. | <i>Path</i> | A trail with distinct vertices. | <i>Connected</i> | A graph where there exists a path between any two vertices. | <i>Component</i> | A maximal connected subgraph. | <i>Tree</i> | A connected acyclic graph. | <i>Free tree</i> | A tree with no root. | <i>DAG</i> | Directed acyclic graph. | <i>Eulerian</i> | Graph with a trail visiting each edge exactly once. | <i>Hamiltonian</i> | Graph with a cycle visiting each vertex exactly once. | <i>Cut</i> | A set of edges whose removal increases the number of components. | <i>Cut-set</i> | A minimal cut. | <i>Cut edge</i> | A size 1 cut. | <i>k-Connected</i> | A graph connected with the removal of any $k-1$ vertices. | <i>k-Tough</i> | $\forall S \subseteq V, S \neq \emptyset$ we have $k \cdot c(G-S) \leq S $. | <i>k-Regular</i> | A graph where all vertices have degree k . | <i>k-Factor</i> | A k -regular spanning subgraph. | <i>Matching</i> | A set of edges, no two of which are adjacent. | <i>Clique</i> | A set of vertices, all of which are adjacent. | <i>Ind. set</i> | A set of vertices, none of which are adjacent. | <i>Vertex cover</i> | A set of vertices which cover all edges. | <i>Planar graph</i> | A graph which can be embedded in the plane. | <i>Plane graph</i> | An embedding of a planar graph. |
| <i>Loop</i> | An edge connecting a vertex to itself. | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | |
| <i>Directed</i> | Each edge has a direction. | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | |
| <i>Simple</i> | Graph with no loops or multi-edges. | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | |
| <i>Walk</i> | A sequence $v_0 e_1 v_1 \dots e_\ell v_\ell$. | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | |
| <i>Trail</i> | A walk with distinct edges. | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | |
| <i>Path</i> | A trail with distinct vertices. | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | |
| <i>Connected</i> | A graph where there exists a path between any two vertices. | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | |
| <i>Component</i> | A maximal connected subgraph. | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | |
| <i>Tree</i> | A connected acyclic graph. | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | |
| <i>Free tree</i> | A tree with no root. | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | |
| <i>DAG</i> | Directed acyclic graph. | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | |
| <i>Eulerian</i> | Graph with a trail visiting each edge exactly once. | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | |
| <i>Hamiltonian</i> | Graph with a cycle visiting each vertex exactly once. | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | |
| <i>Cut</i> | A set of edges whose removal increases the number of components. | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | |
| <i>Cut-set</i> | A minimal cut. | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | |
| <i>Cut edge</i> | A size 1 cut. | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | |
| <i>k-Connected</i> | A graph connected with the removal of any $k-1$ vertices. | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | |
| <i>k-Tough</i> | $\forall S \subseteq V, S \neq \emptyset$ we have $k \cdot c(G-S) \leq S $. | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | |
| <i>k-Regular</i> | A graph where all vertices have degree k . | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | |
| <i>k-Factor</i> | A k -regular spanning subgraph. | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | |
| <i>Matching</i> | A set of edges, no two of which are adjacent. | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | |
| <i>Clique</i> | A set of vertices, all of which are adjacent. | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | |
| <i>Ind. set</i> | A set of vertices, none of which are adjacent. | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | |
| <i>Vertex cover</i> | A set of vertices which cover all edges. | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | |
| <i>Planar graph</i> | A graph which can be embedded in the plane. | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | |
| <i>Plane graph</i> | An embedding of a planar graph. | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | |
| | <p>Notation:</p> <table> <tr> <td>$E(G)$</td><td>Edge set</td></tr> <tr> <td>$V(G)$</td><td>Vertex set</td></tr> <tr> <td>$c(G)$</td><td>Number of components</td></tr> <tr> <td>$G[S]$</td><td>Induced subgraph</td></tr> <tr> <td>$\deg(v)$</td><td>Degree of v</td></tr> <tr> <td>$\Delta(G)$</td><td>Maximum degree</td></tr> <tr> <td>$\delta(G)$</td><td>Minimum degree</td></tr> <tr> <td>$\chi(G)$</td><td>Chromatic number</td></tr> <tr> <td>$\chi_E(G)$</td><td>Edge chromatic number</td></tr> <tr> <td>G^c</td><td>Complement graph</td></tr> <tr> <td>K_n</td><td>Complete graph</td></tr> <tr> <td>K_{n_1, n_2}</td><td>Complete bipartite graph</td></tr> <tr> <td>$r(k, \ell)$</td><td>Ramsey number</td></tr> </table> | $E(G)$ | Edge set | $V(G)$ | Vertex set | $c(G)$ | Number of components | $G[S]$ | Induced subgraph | $\deg(v)$ | Degree of v | $\Delta(G)$ | Maximum degree | $\delta(G)$ | Minimum degree | $\chi(G)$ | Chromatic number | $\chi_E(G)$ | Edge chromatic number | G^c | Complement graph | K_n | Complete graph | K_{n_1, n_2} | Complete bipartite graph | $r(k, \ell)$ | Ramsey number | | | | | | | | | | | | | | | | | | | | | | | | | | |
| $E(G)$ | Edge set | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | |
| $V(G)$ | Vertex set | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | |
| $c(G)$ | Number of components | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | |
| $G[S]$ | Induced subgraph | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | |
| $\deg(v)$ | Degree of v | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | |
| $\Delta(G)$ | Maximum degree | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | |
| $\delta(G)$ | Minimum degree | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | |
| $\chi(G)$ | Chromatic number | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | |
| $\chi_E(G)$ | Edge chromatic number | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | |
| G^c | Complement graph | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | |
| K_n | Complete graph | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | |
| K_{n_1, n_2} | Complete bipartite graph | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | |
| $r(k, \ell)$ | Ramsey number | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | |
| | <p>Geometry</p> <p>Projective coordinates: triples (x, y, z), not all x, y and z zero.</p> $(x, y, z) = (cx, cy, cz) \quad \forall c \neq 0.$ <table> <tr> <td>Cartesian</td><td>Projective</td></tr> <tr> <td>(x, y)</td><td>$(x, y, 1)$</td></tr> <tr> <td>$y = mx + b$</td><td>$(m, -1, b)$</td></tr> <tr> <td>$x = c$</td><td>$(1, 0, -c)$</td></tr> </table> <p>Distance formula, L_p and L_∞ metric:</p> $\sqrt{(x_1 - x_0)^2 + (y_1 - y_0)^2},$ $[x_1 - x_0 ^p + y_1 - y_0 ^p]^{1/p},$ $\lim_{p \rightarrow \infty} [x_1 - x_0 ^p + y_1 - y_0 ^p]^{1/p}.$ <p>Area of triangle (x_0, y_0), (x_1, y_1) and (x_2, y_2):</p> $\frac{1}{2} \operatorname{abs} \begin{vmatrix} x_1 - x_0 & y_1 - y_0 \\ x_2 - x_0 & y_2 - y_0 \end{vmatrix}.$ <p>Angle formed by three points:</p> $\cos \theta = \frac{(x_1, y_1) \cdot (x_2, y_2)}{\ell_1 \ell_2}.$ <p>Line through two points (x_0, y_0) and (x_1, y_1):</p> $\begin{vmatrix} x & y & 1 \\ x_0 & y_0 & 1 \\ x_1 & y_1 & 1 \end{vmatrix} = 0.$ <p>Area of circle, volume of sphere:</p> $A = \pi r^2, \quad V = \frac{4}{3} \pi r^3.$ <p>If I have seen farther than others,</p> | Cartesian | Projective | (x, y) | $(x, y, 1)$ | $y = mx + b$ | $(m, -1, b)$ | $x = c$ | $(1, 0, -c)$ | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | |
| Cartesian | Projective | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | |
| (x, y) | $(x, y, 1)$ | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | |
| $y = mx + b$ | $(m, -1, b)$ | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | |
| $x = c$ | $(1, 0, -c)$ | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | | |

Theoretical Computer Science Cheat Sheet

π

Wallis' identity:

$$\pi = 2 \cdot \frac{2 \cdot 2 \cdot 4 \cdot 4 \cdot 6 \cdot 6 \cdots}{1 \cdot 3 \cdot 3 \cdot 5 \cdot 5 \cdot 7 \cdots}$$

Brouncker's continued fraction expansion:

$$\frac{\pi}{4} = 1 + \cfrac{1^2}{2 + \cfrac{3^2}{2 + \cfrac{5^2}{2 + \cfrac{7^2}{\cdots}}}}$$

Gregory's series:

$$\frac{\pi}{4} = 1 - \frac{1}{3} + \frac{1}{5} - \frac{1}{7} + \frac{1}{9} - \cdots$$

Newton's series:

$$\frac{\pi}{6} = \frac{1}{2} + \frac{1}{2 \cdot 3 \cdot 2^3} + \frac{1 \cdot 3}{2 \cdot 4 \cdot 5 \cdot 2^5} + \cdots$$

Sharp's series:

$$\frac{\pi}{6} = \frac{1}{\sqrt{3}} \left(1 - \frac{1}{3^1 \cdot 3} + \frac{1}{3^2 \cdot 5} - \frac{1}{3^3 \cdot 7} + \cdots \right)$$

Euler's series:

$$\frac{\pi^2}{6} = \frac{1}{1^2} + \frac{1}{2^2} + \frac{1}{3^2} + \frac{1}{4^2} + \frac{1}{5^2} + \cdots$$

$$\frac{\pi^2}{8} = \frac{1}{1^2} + \frac{1}{3^2} + \frac{1}{5^2} + \frac{1}{7^2} + \frac{1}{9^2} + \cdots$$

$$\frac{\pi^2}{12} = \frac{1}{1^2} - \frac{1}{2^2} + \frac{1}{3^2} - \frac{1}{4^2} + \frac{1}{5^2} - \cdots$$

Partial Fractions

Let $N(x)$ and $D(x)$ be polynomial functions of x . We can break down $N(x)/D(x)$ using partial fraction expansion. First, if the degree of N is greater than or equal to the degree of D , divide N by D , obtaining

$$\frac{N(x)}{D(x)} = Q(x) + \frac{N'(x)}{D(x)},$$

where the degree of N' is less than that of D . Second, factor $D(x)$. Use the following rules: For a non-repeated factor:

$$\frac{N(x)}{(x-a)D(x)} = \frac{A}{x-a} + \frac{N'(x)}{D(x)},$$

where

$$A = \left[\frac{N(x)}{D(x)} \right]_{x=a}.$$

For a repeated factor:

$$\frac{N(x)}{(x-a)^m D(x)} = \sum_{k=0}^{m-1} \frac{A_k}{(x-a)^{m-k}} + \frac{N'(x)}{D(x)},$$

where

$$A_k = \frac{1}{k!} \left[\frac{d^k}{dx^k} \left(\frac{N(x)}{D(x)} \right) \right]_{x=a}.$$

The reasonable man adapts himself to the world; the unreasonable persists in trying

Calculus

Derivatives:

1. $\frac{d(cu)}{dx} = c \frac{du}{dx},$
2. $\frac{d(u+v)}{dx} = \frac{du}{dx} + \frac{dv}{dx},$
3. $\frac{d(uv)}{dx} = u \frac{dv}{dx} + v \frac{du}{dx},$
4. $\frac{d(u^n)}{dx} = nu^{n-1} \frac{du}{dx},$
5. $\frac{d(u/v)}{dx} = \frac{v \left(\frac{du}{dx} \right) - u \left(\frac{dv}{dx} \right)}{v^2},$
6. $\frac{d(e^{cu})}{dx} = ce^{cu} \frac{du}{dx},$
7. $\frac{d(c^u)}{dx} = (\ln c)c^u \frac{du}{dx},$
9. $\frac{d(\sin u)}{dx} = \cos u \frac{du}{dx},$
11. $\frac{d(\tan u)}{dx} = \sec^2 u \frac{du}{dx},$
13. $\frac{d(\sec u)}{dx} = \tan u \sec u \frac{du}{dx},$
15. $\frac{d(\arcsin u)}{dx} = \frac{1}{\sqrt{1-u^2}} \frac{du}{dx},$
17. $\frac{d(\arctan u)}{dx} = \frac{1}{1+u^2} \frac{du}{dx},$
19. $\frac{d(\text{arcsec } u)}{dx} = \frac{1}{u\sqrt{1-u^2}} \frac{du}{dx},$
21. $\frac{d(\sinh u)}{dx} = \cosh u \frac{du}{dx},$
23. $\frac{d(\tanh u)}{dx} = \operatorname{sech}^2 u \frac{du}{dx},$
25. $\frac{d(\operatorname{sech } u)}{dx} = -\operatorname{sech } u \tanh u \frac{du}{dx},$
27. $\frac{d(\text{arsinh } u)}{dx} = \frac{1}{\sqrt{1+u^2}} \frac{du}{dx},$
29. $\frac{d(\text{arctanh } u)}{dx} = \frac{1}{1-u^2} \frac{du}{dx},$
31. $\frac{d(\text{arcsech } u)}{dx} = \frac{-1}{u\sqrt{1-u^2}} \frac{du}{dx},$

Integrals:

1. $\int cu \, dx = c \int u \, dx,$
2. $\int (u+v) \, dx = \int u \, dx + \int v \, dx,$
3. $\int x^n \, dx = \frac{1}{n+1} x^{n+1}, \quad n \neq -1,$
4. $\int \frac{1}{x} \, dx = \ln|x|,$
5. $\int e^x \, dx = e^x,$
6. $\int \frac{dx}{1+x^2} = \arctan x,$
8. $\int \sin x \, dx = -\cos x,$
10. $\int \tan x \, dx = -\ln|\cos x|,$
12. $\int \sec x \, dx = \ln|\sec x + \tan x|,$
2. $\int u \, dv = uv - \int v \, du,$
9. $\int \cos x \, dx = \sin x,$
11. $\int \cot x \, dx = \ln|\cos x|,$
13. $\int \csc x \, dx = \ln|\csc x + \cot x|,$

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Calculus Cont.

15. $\int \arccos \frac{x}{a} dx = \arccos \frac{x}{a} - \sqrt{a^2 - x^2}, \quad a > 0,$

16. $\int \arctan \frac{x}{a} dx = x \arctan \frac{x}{a} - \frac{a}{2} \ln(a^2 + x^2),$

17. $\int \sin^2(ax) dx = \frac{1}{2a} (ax - \sin(ax) \cos(ax)),$

18. $\int \cos^2(ax) dx = \frac{1}{2a} (ax + \sin(ax) \cos(ax)),$

19. $\int \sec^2 x dx = \tan x,$

20. $\int \csc^2 x dx = -\cot x,$

21. $\int \sin^n x dx = -\frac{\sin^{n-1} x \cos x}{n} + \frac{n-1}{n} \int \sin^{n-2} x dx,$

22. $\int \cos^n x dx = \frac{\cos^{n-1} x \sin x}{n} + \frac{n-1}{n} \int \cos^{n-2} x dx,$

23. $\int \tan^n x dx = \frac{\tan^{n-1} x}{n-1} - \int \tan^{n-2} x dx, \quad n \neq 1,$

24. $\int \cot^n x dx = -\frac{\cot^{n-1} x}{n-1} - \int \cot^{n-2} x dx, \quad n \neq 1,$

25. $\int \sec^n x dx = \frac{\tan x \sec^{n-1} x}{n-1} + \frac{n-2}{n-1} \int \sec^{n-2} x dx, \quad n \neq 1,$

26. $\int \csc^n x dx = -\frac{\cot x \csc^{n-1} x}{n-1} + \frac{n-2}{n-1} \int \csc^{n-2} x dx, \quad n \neq 1, \quad 27. \int \sinh x dx = \cosh x, \quad 28. \int \cosh x dx = \sinh x,$

29. $\int \tanh x dx = \ln |\cosh x|, \quad 30. \int \coth x dx = \ln |\sinh x|, \quad 31. \int \sech x dx = \arctan \sinh x, \quad 32. \int \csch x dx = \ln |\tanh \frac{x}{2}|,$

33. $\int \sinh^2 x dx = \frac{1}{4} \sinh(2x) - \frac{1}{2}x, \quad 34. \int \cosh^2 x dx = \frac{1}{4} \sinh(2x) + \frac{1}{2}x, \quad 35. \int \sech^2 x dx = \tanh x,$

36. $\int \operatorname{arcsinh} \frac{x}{a} dx = x \operatorname{arcsinh} \frac{x}{a} - \sqrt{x^2 + a^2}, \quad a > 0,$

37. $\int \operatorname{arctanh} \frac{x}{a} dx = x \operatorname{arctanh} \frac{x}{a} + \frac{a}{2} \ln |a^2 - x^2|,$

38. $\int \operatorname{arccosh} \frac{x}{a} dx = \begin{cases} x \operatorname{arccosh} \frac{x}{a} - \sqrt{x^2 + a^2}, & \text{if } \operatorname{arccosh} \frac{x}{a} > 0 \text{ and } a > 0, \\ x \operatorname{arccosh} \frac{x}{a} + \sqrt{x^2 + a^2}, & \text{if } \operatorname{arccosh} \frac{x}{a} < 0 \text{ and } a > 0, \end{cases}$

39. $\int \frac{dx}{\sqrt{a^2 + x^2}} = \ln \left(x + \sqrt{a^2 + x^2} \right), \quad a > 0,$

40. $\int \frac{dx}{a^2 + x^2} = \frac{1}{a} \arctan \frac{x}{a}, \quad a > 0, \quad 41. \int \sqrt{a^2 - x^2} dx = \frac{x}{2} \sqrt{a^2 - x^2} + \frac{a^2}{2} \arcsin \frac{x}{a}, \quad a > 0,$

42. $\int (a^2 - x^2)^{3/2} dx = \frac{x}{8} (5a^2 - 2x^2) \sqrt{a^2 - x^2} + \frac{3a^4}{8} \arcsin \frac{x}{a}, \quad a > 0,$

43. $\int \frac{dx}{\sqrt{a^2 - x^2}} = \arcsin \frac{x}{a}, \quad a > 0, \quad 44. \int \frac{dx}{a^2 - x^2} = \frac{1}{2a} \ln \left| \frac{a+x}{a-x} \right|, \quad 45. \int \frac{dx}{(a^2 - x^2)^{3/2}} = \frac{x}{a^2 \sqrt{a^2 - x^2}},$

46. $\int \sqrt{a^2 \pm x^2} dx = \frac{x}{2} \sqrt{a^2 \pm x^2} \pm \frac{a^2}{2} \ln \left| x + \sqrt{a^2 \pm x^2} \right|,$

47. $\int \frac{dx}{\sqrt{x^2 - a^2}} = \ln \left| x + \sqrt{x^2 - a^2} \right|, \quad a > 0,$

48. $\int \frac{dx}{ax^2 + bx} = \frac{1}{a} \ln \left| \frac{x}{a+bx} \right|,$

49. $\int x \sqrt{a+bx} dx = \frac{2(3bx-2a)(a+bx)^{3/2}}{15b^2},$

50. $\int \frac{\sqrt{a+bx}}{x} dx = 2\sqrt{a+bx} + a \int \frac{1}{x\sqrt{a+bx}} dx,$

51. $\int \frac{x}{\sqrt{a+bx}} dx = \frac{1}{\sqrt{2}} \ln \left| \frac{\sqrt{a+bx} - \sqrt{a}}{\sqrt{a+bx} + \sqrt{a}} \right|, \quad a > 0,$

52. $\int \frac{\sqrt{a^2 - x^2}}{x} dx = \sqrt{a^2 - x^2} - a \ln \left| \frac{a + \sqrt{a^2 - x^2}}{x} \right|,$

53. $\int x \sqrt{a^2 - x^2} dx = -\frac{1}{3} (a^2 - x^2)^{3/2},$

54. $\int x^2 \sqrt{a^2 - x^2} dx = \frac{x}{8} (2x^2 - a^2) \sqrt{a^2 - x^2} + \frac{a^4}{8} \arcsin \frac{x}{a}, \quad a > 0,$

55. $\int \frac{dx}{\sqrt{a^2 - x^2}} = -\frac{1}{a} \ln \left| \frac{a + \sqrt{a^2 - x^2}}{x} \right|,$

56. $\int \frac{x dx}{\sqrt{a^2 - x^2}} = -\sqrt{a^2 - x^2},$

57. $\int \frac{x^2 dx}{\sqrt{a^2 - x^2}} = -\frac{x}{2} \sqrt{a^2 - x^2} + \frac{a^2}{2} \arcsin \frac{x}{a}, \quad a > 0,$

58. $\int \frac{\sqrt{a^2 + x^2}}{x} dx = \sqrt{a^2 + x^2} - a \ln \left| \frac{a + \sqrt{a^2 + x^2}}{x} \right|,$

59. $\int \frac{\sqrt{x^2 - a^2}}{x} dx = \sqrt{x^2 - a^2} - a \arccos \frac{a}{|x|}, \quad a > 0,$

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Calculus Cont.

62. $\int \frac{dx}{x\sqrt{x^2 - a^2}} = \frac{1}{a} \arccos \frac{a}{|x|}, \quad a > 0,$ 63. $\int \frac{dx}{x^2\sqrt{x^2 \pm a^2}} = \mp \frac{\sqrt{x^2 \pm a^2}}{a^2 x},$

64. $\int \frac{x dx}{\sqrt{x^2 \pm a^2}} = \sqrt{x^2 \pm a^2},$ 65. $\int \frac{\sqrt{x^2 \pm a^2}}{x^4} dx = \mp \frac{(x^2 + a^2)^{3/2}}{3a^2 x^3},$

66. $\int \frac{dx}{ax^2 + bx + c} = \begin{cases} \frac{1}{\sqrt{b^2 - 4ac}} \ln \left| \frac{2ax + b - \sqrt{b^2 - 4ac}}{2ax + b + \sqrt{b^2 - 4ac}} \right|, & \text{if } b^2 > 4ac, \\ \frac{2}{\sqrt{4ac - b^2}} \arctan \frac{2ax + b}{\sqrt{4ac - b^2}}, & \text{if } b^2 < 4ac, \end{cases}$

67. $\int \frac{dx}{\sqrt{ax^2 + bx + c}} = \begin{cases} \frac{1}{\sqrt{a}} \ln \left| 2ax + b + 2\sqrt{a}\sqrt{ax^2 + bx + c} \right|, & \text{if } a > 0, \\ \frac{1}{\sqrt{-a}} \arcsin \frac{-2ax - b}{\sqrt{b^2 - 4ac}}, & \text{if } a < 0, \end{cases}$

68. $\int \sqrt{ax^2 + bx + c} dx = \frac{2ax + b}{4a} \sqrt{ax^2 + bx + c} + \frac{4ax - b^2}{8a} \int \frac{dx}{\sqrt{ax^2 + bx + c}},$

69. $\int \frac{x dx}{\sqrt{ax^2 + bx + c}} = \frac{\sqrt{ax^2 + bx + c}}{a} - \frac{b}{2a} \int \frac{dx}{\sqrt{ax^2 + bx + c}},$

70. $\int \frac{dx}{x\sqrt{ax^2 + bx + c}} = \begin{cases} \frac{-1}{\sqrt{c}} \ln \left| \frac{2\sqrt{c}\sqrt{ax^2 + bx + c} + bx + 2c}{x} \right|, & \text{if } c > 0, \\ \frac{1}{\sqrt{-c}} \arcsin \frac{bx + 2c}{|x|\sqrt{b^2 - 4ac}}, & \text{if } c < 0, \end{cases}$

71. $\int x^3 \sqrt{x^2 + a^2} dx = (\frac{1}{3}x^2 - \frac{2}{15}a^2)(x^2 + a^2)^{3/2},$

72. $\int x^n \sin(ax) dx = -\frac{1}{a} x^n \cos(ax) + \frac{n}{a} \int x^{n-1} \cos(ax) dx,$

73. $\int x^n \cos(ax) dx = \frac{1}{a} x^n \sin(ax) - \frac{n}{a} \int x^{n-1} \sin(ax) dx,$

74. $\int x^n e^{ax} dx = \frac{x^n e^{ax}}{a} - \frac{n}{a} \int x^{n-1} e^{ax} dx,$

75. $\int x^n \ln(ax) dx = x^{n+1} \left(\frac{\ln(ax)}{n+1} - \frac{1}{(n+1)^2} \right),$

76. $\int x^n (\ln ax)^m dx = \frac{x^{n+1}}{n+1} (\ln ax)^m - \frac{m}{n+1} \int x^n (\ln ax)^{m-1} dx.$

$$\begin{array}{llll}
 x^1 & x^1 & = & x^{\bar{1}} \\
 x^2 & x^2 + x^1 & = & x^{\bar{2}} - x^{\bar{1}} \\
 x^3 & x^3 + 3x^2 + x^1 & = & x^{\bar{3}} - 3x^{\bar{2}} + x^{\bar{1}} \\
 x^4 & x^4 + 6x^3 + 7x^2 + x^1 & = & x^{\bar{4}} - 6x^{\bar{3}} + 7x^{\bar{2}} - x^{\bar{1}} \\
 x^5 & x^5 + 15x^4 + 25x^3 + 10x^2 + x^1 & = & x^{\bar{5}} - 15x^{\bar{4}} + 25x^{\bar{3}} - 10x^{\bar{2}} + x^{\bar{1}} \\
 x^{\bar{1}} & x^1 & x^{\bar{1}} = & x^1 \\
 x^{\bar{2}} & x^2 + x^1 & x^{\bar{2}} = & x^2 - x^1 \\
 x^{\bar{3}} & x^3 + 3x^2 + 2x^1 & x^{\bar{3}} = & x^3 - 3x^2 + 2x^1 \\
 \bar{x} & & &
 \end{array}$$

Finite Calculus

Difference, shift operators:

$$\Delta f(x) = f(x+1) - f(x),$$

$$\text{E } f(x) = f(x+1).$$

Fundamental Theorem:

$$f(x) = \Delta F(x) \Leftrightarrow \sum_b f(x) \delta x = F(x) + C.$$

$$\sum_a f(x) \delta x = \sum_{i=a}^{b-1} f(i).$$

Differences:

$$\Delta(cu) = c\Delta u, \quad \Delta(u+v) = \Delta u + \Delta v,$$

$$\Delta(uv) = u\Delta v + \text{E } v \Delta u,$$

$$\Delta(x^n) = nx^{n-1},$$

$$\Delta(H_x) = x^{-1}, \quad \Delta(2^x) = 2^x,$$

$$\Delta(c^x) = (c-1)c^x, \quad \Delta(\binom{x}{m}) = \binom{x}{m-1}.$$

Sums:

$$\sum cu \delta x = c \sum u \delta x,$$

$$\sum (u+v) \delta x = \sum u \delta x + \sum v \delta x,$$

$$\sum u \Delta v \delta x = uv - \sum \text{E } v \Delta u \delta x,$$

$$\sum x^n \delta x = \frac{x^{n+1}}{n+1}, \quad \sum x^{-1} \delta x = H_x,$$

$$\sum c^x \delta x = \frac{c^x}{c-1}, \quad \sum \binom{x}{m} \delta x = \binom{x}{m+1}.$$

Falling Factorial Powers:

$$x^{\underline{n}} = x(x-1) \cdots (x-n+1), \quad n > 0,$$

$$x^{\underline{0}} = 1,$$

$$x^{\overline{n}} = \frac{1}{(x+1) \cdots (x+n)}, \quad n < 0,$$

$$x^{\underline{n+m}} = x^{\underline{m}} (x-m)^{\underline{n}}.$$

Rising Factorial Powers:

$$x^{\overline{n}} = x(x+1) \cdots (x+n-1), \quad n > 0,$$

$$x^{\overline{0}} = 1,$$

$$x^{\overline{n}} = \frac{1}{(x-1) \cdots (x-|n|)}, \quad n < 0,$$

$$x^{\overline{n+m}} = x^{\overline{m}} (x+m)^{\overline{n}}.$$

Conversion:

$$x^{\underline{n}} = (-1)^n (-x)^{\overline{n}} = (x-n+1)^{\overline{n}}$$

$$= 1/(x+1)^{\overline{-n}},$$

$$x^{\overline{n}} = (-1)^n (-x)^{\underline{n}} = (x+n-1)^{\underline{n}}$$

$$= 1/(x-1)^{\underline{-n}},$$

$$x^{\underline{n}} = \sum_{k=1}^n \binom{n}{k} x^k = \sum_{k=1}^n \binom{n}{k} (-1)^{n-k} x^{\overline{k}},$$

$$x^{\overline{n}} = \sum_{k=1}^n \binom{n}{k} (-1)^{n-k} x^k,$$

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Series

Taylor's series:

$$f(x) = f(a) + (x-a)f'(a) + \frac{(x-a)^2}{2}f''(a) + \dots = \sum_{i=0}^{\infty} \frac{(x-a)^i}{i!} f^{(i)}(a).$$

Expansions:

| | | |
|--|---|--|
| $\frac{1}{1-x}$ | $= 1 + x + x^2 + x^3 + x^4 + \dots$ | $= \sum_{i=0}^{\infty} x^i,$ |
| $\frac{1}{1-cx}$ | $= 1 + cx + c^2x^2 + c^3x^3 + \dots$ | $= \sum_{i=0}^{\infty} c^i x^i,$ |
| $\frac{1}{1-x^n}$ | $= 1 + x^n + x^{2n} + x^{3n} + \dots$ | $= \sum_{i=0}^{\infty} x^{ni},$ |
| $\frac{x}{(1-x)^2}$ | $= x + 2x^2 + 3x^3 + 4x^4 + \dots$ | $= \sum_{i=0}^{\infty} ix^i,$ |
| $x^k \frac{d^n}{dx^n} \left(\frac{1}{1-x} \right)$ | $= x + 2^n x^2 + 3^n x^3 + 4^n x^4 + \dots$ | $= \sum_{i=0}^{\infty} i^n x^i,$ |
| e^x | $= 1 + x + \frac{1}{2}x^2 + \frac{1}{6}x^3 + \dots$ | $= \sum_{i=0}^{\infty} \frac{x^i}{i!},$ |
| $\ln(1+x)$ | $= x - \frac{1}{2}x^2 + \frac{1}{3}x^3 - \frac{1}{4}x^4 - \dots$ | $= \sum_{i=1}^{\infty} (-1)^{i+1} \frac{x^i}{i},$ |
| $\ln \frac{1}{1-x}$ | $= x + \frac{1}{2}x^2 + \frac{1}{3}x^3 + \frac{1}{4}x^4 + \dots$ | $= \sum_{i=1}^{\infty} \frac{x^i}{i},$ |
| $\sin x$ | $= x - \frac{1}{3!}x^3 + \frac{1}{5!}x^5 - \frac{1}{7!}x^7 + \dots$ | $= \sum_{i=0}^{\infty} (-1)^i \frac{x^{2i+1}}{(2i+1)!},$ |
| $\cos x$ | $= 1 - \frac{1}{2!}x^2 + \frac{1}{4!}x^4 - \frac{1}{6!}x^6 + \dots$ | $= \sum_{i=0}^{\infty} (-1)^i \frac{x^{2i}}{(2i)!},$ |
| $\tan^{-1} x$ | $= x - \frac{1}{3}x^3 + \frac{1}{5}x^5 - \frac{1}{7}x^7 + \dots$ | $= \sum_{i=0}^{\infty} (-1)^i \frac{x^{2i+1}}{(2i+1)},$ |
| $(1+x)^n$ | $= 1 + nx + \frac{n(n-1)}{2}x^2 + \dots$ | $= \sum_{i=0}^{\infty} \binom{n}{i} x^i,$ |
| $\frac{1}{(1-x)^{n+1}}$ | $= 1 + (n+1)x + \binom{n+2}{2}x^2 + \dots$ | $= \sum_{i=0}^{\infty} \binom{i+n}{i} x^i,$ |
| $\frac{x}{e^x - 1}$ | $= 1 - \frac{1}{2}x + \frac{1}{12}x^2 - \frac{1}{720}x^4 + \dots$ | $= \sum_{i=0}^{\infty} \frac{B_i x^i}{i!},$ |
| $\frac{1}{2x}(1 - \sqrt{1-4x})$ | $= 1 + x + 2x^2 + 5x^3 + \dots$ | $= \sum_{i=0}^{\infty} \frac{1}{i+1} \binom{2i}{i} x^i,$ |
| $\frac{1}{\sqrt{1-4x}}$ | $= 1 + x + 2x^2 + 6x^3 + \dots$ | $= \sum_{i=0}^{\infty} \binom{2i}{i} x^i,$ |
| $\frac{1}{\sqrt[3]{1-4x}} \left(\frac{1-\sqrt{1-4x}}{2x} \right)^n$ | $= 1 + (2+n)x + \binom{4+n}{2}x^2 + \dots$ | $= \sum_{i=0}^{\infty} \binom{2i+n}{i} x^i,$ |
| $\frac{1}{1-x} \ln \frac{1}{1-x}$ | $= x + \frac{3}{2}x^2 + \frac{11}{6}x^3 + \frac{25}{12}x^4 + \dots$ | $= \sum_{i=1}^{\infty} H_i x^i,$ |
| $\frac{1}{2} \left(\ln \frac{1}{1-x} \right)^2$ | $= \frac{1}{2}x^2 + \frac{3}{4}x^3 + \frac{11}{24}x^4 + \dots$ | $= \sum_{i=2}^{\infty} \frac{H_{i-1} x^i}{i},$ |
| $\frac{x}{1 - \frac{x}{x^2}}$ | $= x + x^2 + 2x^3 + 3x^4 + \dots$ | $= \sum_{i=1}^{\infty} F_i x^i,$ |

Ordinary power series:

$$A(x) = \sum_{i=0}^{\infty} a_i x^i.$$

Exponential power series:

$$A(x) = \sum_{i=0}^{\infty} a_i \frac{x^i}{i!}.$$

Dirichlet power series:

$$A(x) = \sum_{i=1}^{\infty} \frac{a_i}{i^x}.$$

Binomial theorem:

$$(x+y)^n = \sum_{k=0}^n \binom{n}{k} x^{n-k} y^k.$$

Difference of like powers:

$$x^n - y^n = (x-y) \sum_{k=0}^{n-1} x^{n-1-k} y^k.$$

For ordinary power series:

$$\alpha A(x) + \beta B(x) = \sum_{i=0}^{\infty} (\alpha a_i + \beta b_i) x^i,$$

$$x^k A(x) = \sum_{i=k}^{\infty} a_{i-k} x^i,$$

$$\frac{A(x) - \sum_{i=0}^{k-1} a_i x^i}{x^k} = \sum_{i=0}^{\infty} a_{i+k} x^i,$$

$$A(cx) = \sum_{i=0}^{\infty} c^i a_i x^i,$$

$$A'(x) = \sum_{i=0}^{\infty} (i+1) a_{i+1} x^i,$$

$$xA'(x) = \sum_{i=1}^{\infty} ia_i x^i,$$

$$\int A(x) dx = \sum_{i=1}^{\infty} \frac{a_{i-1}}{i} x^i,$$

$$\frac{A(x) + A(-x)}{2} = \sum_{i=0}^{\infty} a_{2i} x^{2i},$$

$$\frac{A(x) - A(-x)}{2} = \sum_{i=0}^{\infty} a_{2i+1} x^{2i+1}.$$

Summation: If $b_i = \sum_{j=0}^i a_j$ then

$$B(x) = \frac{1}{1-x} A(x).$$

Convolution:

$$A(x)B(x) = \sum_{i=0}^{\infty} \left(\sum_{j=0}^i a_j b_{i-j} \right) x^i.$$

Theoretical Computer Science Cheat Sheet

Series

Expansions:

$$\begin{aligned}
 \frac{1}{(1-x)^{n+1}} \ln \frac{1}{1-x} &= \sum_{i=0}^{\infty} (H_{n+i} - H_n) \binom{n+i}{i} x^i, \\
 x^{\bar{n}} &= \sum_{i=0}^{\infty} \binom{n}{i} x^i, \\
 \left(\ln \frac{1}{1-x}\right)^n &= \sum_{i=0}^{\infty} \binom{i}{n} \frac{n!x^i}{i!}, \\
 \tan x &= \sum_{i=1}^{\infty} (-1)^{i-1} \frac{2^{2i}(2^{2i}-1)B_{2i}x^{2i-1}}{(2i)!}, \\
 \frac{1}{\zeta(x)} &= \sum_{i=1}^{\infty} \frac{\mu(i)}{i^x}, \\
 \zeta(x) &= \prod_p \frac{1}{1-p^{-x}}, \\
 \zeta^2(x) &= \sum_{i=1}^{\infty} \frac{d(i)}{x^i} \quad \text{where } d(n) = \sum_{d|n} 1, \\
 \zeta(x)\zeta(x-1) &= \sum_{i=1}^{\infty} \frac{S(i)}{x^i} \quad \text{where } S(n) = \sum_{d|n} d, \\
 \zeta(2n) &= \frac{2^{2n-1}|B_{2n}|}{(2n)!} \pi^{2n}, \quad n \in \mathbb{N}, \\
 \frac{x}{\sin x} &= \sum_{i=0}^{\infty} (-1)^{i-1} \frac{(4^i-2)B_{2i}x^{2i}}{(2i)!}, \\
 \left(\frac{1-\sqrt{1-4x}}{2x}\right)^n &= \sum_{i=0}^{\infty} \frac{n(2i+n-1)!}{i!(n+i)!} x^i, \\
 e^x \sin x &= \sum_{i=1}^{\infty} \frac{2^{i/2} \sin \frac{i\pi}{4}}{i!} x^i, \\
 \sqrt{\frac{1-\sqrt{1-x}}{x}} &= \sum_{i=0}^{\infty} \frac{(4i)!}{16^i \sqrt{2}(2i)!(2i+1)!} x^i, \\
 \left(\frac{\arcsin x}{x}\right)^2 &= \sum_{i=0}^{\infty} \frac{4^i i!^2}{(i+1)(2i+1)!} x^{2i}.
 \end{aligned}$$

Cramer's Rule

If we have equations:

$$a_{1,1}x_1 + a_{1,2}x_2 + \cdots + a_{1,n}x_n = b_1$$

$$a_{2,1}x_1 + a_{2,2}x_2 + \cdots + a_{2,n}x_n = b_2$$

$$\vdots \quad \vdots \quad \vdots$$

$$a_{n,1}x_1 + a_{n,2}x_2 + \cdots + a_{n,n}x_n = b_n$$

Let $A = (a_{i,j})$ and B be the column matrix (b_i) . Then there is a unique solution iff $\det A \neq 0$. Let A_i be A with column i replaced by B . Then

$$x_i = \frac{\det A_i}{\det A}.$$

Improvement makes strait roads, but the crooked

$$\begin{aligned}
 \left(\frac{1}{x}\right)^{-n} &= \sum_{i=0}^{\infty} \binom{i}{n} x^i, \\
 (e^x - 1)^n &= \sum_{i=0}^{\infty} \binom{i}{n} \frac{n!x^i}{i!}, \\
 x \cot x &= \sum_{i=0}^{\infty} \frac{(-4)^i B_{2i} x^{2i}}{(2i)!}, \\
 \zeta(x) &= \sum_{i=1}^{\infty} \frac{1}{i^x}, \\
 \frac{\zeta(x-1)}{\zeta(x)} &= \sum_{i=1}^{\infty} \frac{\phi(i)}{i^x},
 \end{aligned}$$

Escher



Stieltjes Integration

If G is continuous in the interval $[a, b]$ and F is nondecreasing then

$$\int_a^b G(x) dF(x)$$

exists. If $a \leq b \leq c$ then

$$\int_a^c G(x) dF(x) = \int_a^b G(x) dF(x) + \int_b^c G(x) dF(x).$$

If the integrals involved exist

$$\begin{aligned}
 \int_a^b (G(x) + H(x)) dF(x) &= \int_a^b G(x) dF(x) + \int_a^b H(x) dF(x), \\
 \int_a^b G(x) d(F(x) + H(x)) &= \int_a^b G(x) dF(x) + \int_a^b G(x) dH(x), \\
 \int_a^b c \cdot G(x) dF(x) &= \int_a^b G(x) d(c \cdot F(x)) = c \int_a^b G(x) dF(x), \\
 \int_a^b G(x) dF(x) &= G(b)F(b) - G(a)F(a) - \int_a^b F(x) dG(x).
 \end{aligned}$$

If the integrals involved exist, and F possesses a derivative F' at every point in $[a, b]$ then

$$\int_a^b G(x) dF(x) = \int_a^b G(x) F'(x) dx.$$

Fibonacci Numbers

$$0, 1, 2, 3, 5, 8, 13, 21, 34, 55, 89, \dots$$

Definitions:

$$F_i = F_{i-1} + F_{i-2}, \quad F_0 = F_1 = 1,$$

$$F_{-i} = (-1)^{i-1} F_i,$$

$$F_i = \frac{1}{\sqrt{5}} (\phi^i - \hat{\phi}^i),$$

Cassini's identity: for $i > 0$:

$$F_{i+1}F_{i-1} - F_i^2 = (-1)^i.$$

Additive rule:

$$F_{n+k} = F_k F_{n+1} + F_{k-1} F_n,$$

$$F_{2n} = F_n F_{n+1} + F_{n-1} F_n.$$

Calculation by matrices:

| | | | | | | | | | |
|----|----|----|----|----|----|----|----|----|----|
| 00 | 47 | 18 | 76 | 29 | 93 | 85 | 34 | 61 | 52 |
| 86 | 11 | 57 | 28 | 70 | 39 | 94 | 45 | 02 | 63 |
| 95 | 80 | 22 | 67 | 38 | 71 | 49 | 56 | 13 | 04 |
| 59 | 96 | 81 | 33 | 07 | 48 | 72 | 60 | 24 | 15 |
| 73 | 69 | 90 | 82 | 44 | 17 | 58 | 01 | 35 | 26 |
| 68 | 74 | 09 | 91 | 83 | 55 | 27 | 12 | 46 | 30 |
| 37 | 08 | 75 | 19 | 92 | 84 | 66 | 23 | 50 | 41 |
| 14 | 25 | 36 | 40 | 51 | 62 | 03 | 77 | 88 | 99 |
| 21 | 32 | 43 | 54 | 65 | 06 | 10 | 89 | 97 | 78 |
| 42 | 53 | 64 | 05 | 16 | 20 | 31 | 98 | 79 | 87 |

The Fibonacci number system:
Every integer n has a unique representation

$$n = F_{k_1} + F_{k_2} + \cdots + F_{k_m}$$